



# Efficient Schwarz Preconditioning Techniques for Nonlinear Problems Using FROSch

Alexander Heinlein<sup>1</sup>

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Based on joint work with Axel Klawonn and Lea Saßmannshausen (University of Cologne) and Mauro Perego, Sivasankaran Rajamanickam, and Ichitaro Yamazaki (Sandia National Laboratories)

### Solving A Model Problem



Consider a diffusion model problem:

$$-\nabla \cdot (\alpha(x)\nabla u(x)) = f \quad \text{in } \Omega = [0,1]^2,$$
$$u = 0 \quad \text{on } \partial\Omega.$$

Discretization using finite elements yields a **sparse** linear system of equations

$$Ku = f$$
.

#### **Direct solvers**

For fine meshes, solving the system using a direct solver is not feasible due to **superlinear complexity and memory cost**.

#### **Iterative solvers**

Iterative solvers are efficient for solving sparse linear systems of equations, however, the convergence rate generally depends on the condition number  $\kappa$  (*A*). It deteriorates, e.g., for

- fine meshes, that is, small element sizes h

### Solving A Model Problem



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- fine meshes, that is, small element sizes h
- large contrasts  $\frac{\max_{x} \alpha(x)}{\min_{x} \alpha(x)}$

 $\Rightarrow$  We introduce a preconditioner  $M^{-1} \approx A^{-1}$  to improve the condition number:

$$M^{-1}Au = M^{-1}f$$

### **Two-Level Schwarz Preconditioners**

**One-level Schwarz preconditioner** 





Based on an overlapping domain decomposition, we define a one-level Schwarz operator

$$\boldsymbol{M}_{\text{OS-1}}^{-1}\boldsymbol{K} = \sum_{i=1}^{N} \boldsymbol{R}_{i}^{T}\boldsymbol{K}_{i}^{-1}\boldsymbol{R}_{i}\boldsymbol{K}_{i}$$

where  $\mathbf{R}_i$  and  $\mathbf{R}_i^{\mathsf{T}}$  are restriction and prolongation operators corresponding to  $\Omega'_i$ , and  $\mathbf{K}_i := \mathbf{R}_i \mathbf{K} \mathbf{R}_i^{\mathsf{T}}$ .

Condition number estimate:

$$\kappa\left( \pmb{M}_{\mathsf{OS-1}}^{-1}\pmb{K} 
ight) \leq C\left(1+rac{1}{H\delta}
ight)$$

with subdomain size H and overlap width  $\delta$ .

### Lagrangian coarse space





The two-level overlapping Schwarz operator reads

$$\boldsymbol{M}_{\text{OS-2}}^{-1}\boldsymbol{K} = \underbrace{\boldsymbol{\Phi}\boldsymbol{K}_{0}^{-1}\boldsymbol{\Phi}^{\mathsf{T}}\boldsymbol{K}}_{\text{coarse level - global}} + \underbrace{\sum_{i=1}^{N}\boldsymbol{R}_{i}^{\mathsf{T}}\boldsymbol{K}_{i}^{-1}\boldsymbol{R}_{i}\boldsymbol{K}}_{\text{first level - local}},$$

where  $\Phi$  contains the coarse basis functions and  $K_0 := \Phi^T K \Phi$ ; cf., e.g., Toselli, Widlund (2005). The construction of a Lagrangian coarse basis requires a coarse triangulation.

Condition number estimate:

$$\kappa\left(\boldsymbol{M}_{\mathsf{OS-2}}^{-1}\boldsymbol{K}\right) \leq C\left(1+rac{\boldsymbol{H}}{\delta}
ight)$$

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### **Two-Level Schwarz Preconditioners**



## FROSch (Fast and Robust Overlapping Schwarz) Framework in Trilinos





### Software

- Object-oriented C++ domain decomposition solver framework with MPI-based distributed memory parallelization
- Part of TRILINOS with support for both parallel linear algebra packages EPETRA and TPETRA
- Node-level parallelization and performance portability on CPU and GPU architectures through KOKKOS and KOKKOSKERNELS
- Accessible through unified  $\mathrm{TRILINOS}$  solver interface  $\mathrm{STRATIMIKOS}$

### Methodology

- Parallel scalable multi-level Schwarz domain decomposition
   preconditioners
- · Algebraic construction based on the parallel distributed system matrix
- Extension-based coarse spaces

### Team (active)

- Alexander Heinlein (TU Delft)
- Siva Rajamanickam (Sandia)
- Friederike Röver (TUBAF)
- Ichitaro Yamazaki (Sandia)

- Axel Klawonn (Uni Cologne)
- Oliver Rheinbach (TUBAF)
- Lea Saßmannshausen (Uni Cologne)

### Overlapping domain decomposition

In FROSch, the overlapping subdomains  $\Omega'_1, ..., \Omega'_N$  are constructed by **recursively adding layers of elements** to the nonoverlapping subdomains; this can be performed based on the sparsity pattern of K.



Nonoverlapping DD

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 $\text{Overlap } \delta = 1h$ 



Overlap  $\delta = 2h$ 

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Nonoverlapping DD



Overlap  $\delta = 1h$ 



Overlap  $\delta = 2h$ 

### Computation of the overlapping matrices

The overlapping matrices

$$oldsymbol{K}_i = oldsymbol{R}_i oldsymbol{K} oldsymbol{R}_i^T$$

can easily be extracted from K since  $R_i$  is just a global-to-local index mapping.

1. Identification interface components



Identification from parallel distribution of matrix:







### 3. Interface basis

 null space basis

 (e.g., linear elasticity: translations, linearized rotation(s))

 ×

The interface values of the basis of the coarse space is obtained by **multiplication with the null space**.

### 2. Interface partition of unity (IPOU)

vertex & edge functions





Based on the interface components, construct an **interface partition of** unity:

$$\sum_i \pi_i = 1$$
 on  $\Gamma$ 



### 4. Extension into the interior

edge basis function

vertex basis function



The values in the interior of the subdomains are computed via the **extension operator**:

$$\Phi = \begin{bmatrix} \Phi_I \\ \Phi_{\Gamma} \end{bmatrix} = \begin{bmatrix} -K_{II}^{-1}K_{\Gamma I}^{T}\Phi_{\Gamma} \\ \Phi_{\Gamma} \end{bmatrix}.$$

(For elliptic problems: energy-minimizing extension)







### 2. Interface partition of unity (IPOU)

vertex & edge functions



vertex functions

Based on the interface components, construct an interface partition of unity:

$$\sum\nolimits_{i} \pi_{i} = 1 ext{ on } \Gamma$$









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1. Identification interface components
 K = 
 Identification from parallel distribution of matrix:
 distributed map
 overlapping map
 repeated map
 interface components
 3. Interface basis



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### Examples of FROSch Coarse Spaces

### GDSW (Generalized Dryja-Smith-Widlund)





- Dohrmann, Klawonn, Widlund (2008)
- Dohrmann, Widlund (2009, 2010, 2012)

### MsFEM (Multiscale Finite Element Method)





- Hou (1997), Efendiev and Hou (2009)
- Buck, Iliev, and Andrä (2013)
- H., Klawonn, Knepper, Rheinbach (2018)

### **RGDSW** (Reduced dimension GDSW)





- Dohrmann, Widlund (2017)
- H., Klawonn, Knepper, Rheinbach, Widlund (2022)

### Q1 Lagrangian / piecewise bilinear





**Piecewise linear** interface partition of unity functions and a **structured domain decomposition**.

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### Monolithic (R)GDSW Preconditioners for CFD Simulations

Consider the discrete saddle point problem

$$\mathcal{A}_{X} = \begin{bmatrix} \mathbf{K} & \mathbf{B}^{\top} \\ \mathbf{B} & \mathbf{0} \end{bmatrix} \begin{bmatrix} \mathbf{u} \\ \mathbf{p} \end{bmatrix} = \begin{bmatrix} \mathbf{f} \\ \mathbf{0} \end{bmatrix} = \mathbf{6}.$$

### Monolithic GDSW preconditioner

We construct a monolithic GDSW preconditioner

$$\mathcal{M}_{\mathsf{GDSW}}^{-1} = \phi \mathcal{R}_0^{-1} \phi^\top + \sum_{i=1}^N \mathcal{R}_i^\top \mathcal{R}_i^{-1} \mathcal{R}_i,$$

with block matrices  $\mathcal{A}_0 = \phi^\top \mathcal{A} \phi$ ,  $\mathcal{A}_i = \mathcal{R}_i \mathcal{A} \mathcal{R}_i^\top$ , and

$$\mathcal{R}_i = \begin{bmatrix} \mathcal{R}_{u,i} & \mathbf{0} \\ \mathbf{0} & \mathcal{R}_{p,i} \end{bmatrix} \quad \text{and} \quad \phi = \begin{bmatrix} \Phi_{u,u_0} & \Phi_{u,p_0} \\ \Phi_{p,u_0} & \Phi_{p,p_0} \end{bmatrix}.$$

Using  $\mathcal{A}$  to compute extensions:  $\phi_I = -\mathcal{A}_{II}^{-1}\mathcal{A}_{I\Gamma}\phi_{\Gamma}$ ; cf. Heinlein, Hochmuth, Klawonn (2019, 2020).







Stokes flow

Navier-Stokes flow

### **Related work:**

- Original work on monolithic Schwarz preconditioners: Klawonn and Pavarino (1998, 2000)
- Other publications on monolithic Schwarz preconditioners: e.g., Hwang and Cai (2006), Barker and Cai (2010), Wu and Cai (2014), and the presentation Dohrmann (2010) at the Workshop on Adaptive Finite Elements and Domain Decomposition Methods in Milan.

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### Monolithic (R)GDSW Preconditioners for CFD Simulations

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#### Block diagonal & triangular preconditioners

Block-diagonal preconditioner:

$$\boldsymbol{M}_{\mathrm{D}}^{-1} = \begin{bmatrix} \boldsymbol{K}^{-1} & \boldsymbol{0} \\ \boldsymbol{0} & \boldsymbol{S}^{-1} \end{bmatrix} \approx \begin{bmatrix} \boldsymbol{M}_{\mathrm{GDSW}}^{-1}(\boldsymbol{K}) & \boldsymbol{0} \\ \boldsymbol{0} & \boldsymbol{M}_{\mathrm{OS-1}}^{-1}(\boldsymbol{M}_{\boldsymbol{P}}) \end{bmatrix}$$

Block-triangular preconditioner:

$$\begin{split} m_{\mathsf{T}}^{-1} &= \begin{bmatrix} \boldsymbol{K}^{-1} & \boldsymbol{0} \\ -\boldsymbol{S}^{-1}\boldsymbol{B}\boldsymbol{K}^{-1} & \boldsymbol{S}^{-1} \end{bmatrix} \\ &\approx \begin{bmatrix} \boldsymbol{M}_{\mathsf{GDSW}}^{-1}(\boldsymbol{K}) & \boldsymbol{0} \\ -\boldsymbol{M}_{\mathsf{OS}-1}^{-1}(\boldsymbol{M}_{\boldsymbol{\rho}})\boldsymbol{B}\boldsymbol{M}_{\mathsf{GDSW}}^{-1}(\boldsymbol{K}) & \boldsymbol{M}_{\mathsf{OS}-1}^{-1}(\boldsymbol{M}_{\boldsymbol{\rho}}) \end{bmatrix} \end{split}$$

### Monolithic vs. block prec. (Stokes)



prec.	# MPI ranks	64	256	1024	4 096
mono.	time	154.7 s	170.0 s	175.8 s	188.7 s
	effic.	100 %	91 %	88 %	82 %
tuiona	time	309.4 s	329.1 s	359.8 s	396.7 s
triang.	effic.	50 %	47 %	43 %	39 %
1.	time	736.7 s	859.4 s	966.9 s	$1105.0\mathrm{s}$
diag.	effic.	21 %	18%	16%	14 %

Computations performed on magnitUDE (University Duisburg-Essen).

### Monolithic (R)GDSW Preconditioners for CFD Simulations

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with block matrices  $\mathcal{A}_0 = \phi^\top \mathcal{A} \phi$ ,  $\mathcal{A}_i = \mathcal{R}_i \mathcal{A} \mathcal{R}_i^\top$ .

### SIMPLE block preconditioner

We employ the SIMPLE (Semi-Implicit Method for Pressure Linked Equations) block preconditioner

$$m_{\mathsf{SIMPLE}}^{-1} = \begin{bmatrix} \mathbf{I} & -\mathbf{D}^{-1}\mathbf{B} \\ \mathbf{0} & \alpha \mathbf{I} \end{bmatrix} \begin{bmatrix} \mathbf{K}^{-1} & \mathbf{0} \\ -\hat{\mathbf{S}}^{-1}\mathbf{B}\mathbf{K}^{-1} & \hat{\mathbf{S}}^{-1} \end{bmatrix};$$

see Patankar and Spalding (1972). Here,

- $\hat{\boldsymbol{S}} = -\boldsymbol{B}\boldsymbol{D}^{-1}\boldsymbol{B}^{\top}$ , with  $\boldsymbol{D} = \operatorname{diag} \boldsymbol{K}$
- $\alpha$  is an under-relaxation parameter

We approximate the inverses using (R)GDSW preconditioners.

### Monolithic vs. SIMPLE preconditioner



#### Steady-state Navier-Stokes equations

prec.	# MPI ranks	243	1 1 2 5	15 562
Monolithic	setup	39.6 s	57.9 s	95.5 s
RGDSW	solve	57.6 s	69.2 s	74.9 s
(FROSCH)	total	97.2 s	127.7 s	170.4 s
(FROSCH) SIMPLE	total setup	<b>97.2 s</b> 39.2 s	<b>127.7 s</b> 38.2 s	<b>170.4 s</b> 68.6 s
(FROSCH) SIMPLE RGDSW (Теко	total setup solve	<b>97.2 s</b> 39.2 s <b>86.2 s</b>	<b>127.7 s</b> 38.2 s <b>106.6 s</b>	<b>170.4 s</b> 68.6 s <b>127.4 s</b>

Computations on Piz Daint (CSCS). Implementation in the finite element software FEDDLib.

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### Monolithic Vs. Block (R)GDSW Preconditioners for CFD Simulations



Problem: Steady-state Navier–Stokes equations

Computations on Fritz (FAU). Implementation in the finite element software FEDDLIB.

### Monolithic GDSW preconditioner

$$\mathcal{M}_{\mathsf{GDSW}}^{-1} = \phi \mathcal{R}_0^{-1} \phi^\top + \sum_{i=1}^N \mathcal{R}_i^\top \mathcal{R}_i^{-1} \mathcal{R}_i$$

### SIMPLE block preconditioner

$$\boldsymbol{m}_{\mathsf{SIMPLE}}^{-1} = \begin{bmatrix} \boldsymbol{I} & -\boldsymbol{D}^{-1}\boldsymbol{B} \\ \boldsymbol{0} & \alpha \boldsymbol{I} \end{bmatrix} \begin{bmatrix} \boldsymbol{K}^{-1} & \boldsymbol{0} \\ -\hat{\boldsymbol{S}}^{-1}\boldsymbol{B}\boldsymbol{K}^{-1} & \hat{\boldsymbol{S}}^{-1} \end{bmatrix}$$

### **P1–P1 stabilized,** H/h = 20

prec.	# MPI ranks	243	1 1 2 5	4 608
Monolithic	# its.	57.8(5)	71.6(5)	79.4(5)
	setup	39.6 s	50.9 s	49.8 s
(EDOCarr)	solve	38.2 s	58.7 s	71.7 s
(FROSCH)	total	77.8 s	109.8 s	121.5 s
SIMPLE	# its.	168.4(5)	196.8(5)	200.0(5)
RGDSW	setup	21.2 s	32.2 s	26.9 s
(Теко &	solve	106.2 s	156.0 s	175.0 s
FROSCH)	total	127.4 s	188.2 s	201.9 s

**P2–P1,** *H*/*h* = 9

prec.	# MPI ranks	243	1 1 2 5	4 608
Monolithic	# its.	84.2(6)	100.4(5)	108.6(5)
	setup	44.2 s	48.5 s	49.7 s
(EDOCarr)	solve	50.0 s	63.9 s	88.0 s
(FROSCH)	total	94.2 s	112.4 s	137.7 s
SIMPLE	$\overline{\#}$ its.	157.5(6)	161.8(5)	169.8(5)
RGDSW	setup	26.8 s	31.7 s	28.5 s
(Теко &	solve	84.8 s	90.4 s	111.5 s
FROSCH)	total	111.6 s	122.1 s	140.0 s

Heinlein, Klawonn, and Saßmannshausen (in preparation)

### **FROSch Preconditioners for Land Ice Simulations**

### Stationary velocity problem

We use a first-order (or Blatter-Pattyn) approximation of the Stokes equations

 $\begin{cases} -\nabla \cdot (2\mu \,\dot{\epsilon}_1) &= -\rho_i \, |\boldsymbol{g}| \, \partial_x s, \\ -\nabla \cdot (2\mu \,\dot{\epsilon}_2) &= -\rho_i \, |\boldsymbol{g}| \, \partial_y s, \end{cases}$ 

with ice density  $\rho_i$ , ice surface elevation *s*, gravity acceleration *g*, and strain rates  $\dot{\epsilon}_1$ and  $\dot{\epsilon}_2$ ; cf. Blatter (1995) and Pattyn (2003).

Ice viscosity modeled by Glen's law:  $\mu = \frac{1}{2}A(T)^{-\frac{1}{n}}\dot{\epsilon_e}^{\frac{1-n}{n}}$ .

### Stationary temperature problem

The enthalpy equation reads

$$abla \cdot \boldsymbol{q}(h) + \boldsymbol{u} \cdot \nabla h = 4\mu \,\epsilon_e^2$$

with the enthalpy flux

$$q(h) = \left\{ egin{array}{c} rac{k}{
ho_i c_i} 
abla h, & \quad ext{for cold ice } (h \leq h_m), \ rac{k}{
ho_i c_i} 
abla h_m + 
ho_w L oldsymbol{j}(h), & \quad ext{for temperate ice.} \end{array} 
ight.$$

**Set of complex boundary conditions**: Dirichlet, Neumann, Robin, and Stefan boundary and coupling conditions.



### **FROSch** preconditioners



We compute the **nonoverlapping domain decomposition** based on the **surface mesh**.

For the coupled problem, we construct a **monolithic** 

two-level (R)GDSW preconditioner (Heinlein, Hochmuth, Klawonn (2019, 2020))

$$\mathcal{M}_{\mathrm{GDSW}}^{-1} = \phi \mathcal{R}_0^{-1} \phi^{\mathsf{T}} + \sum_{i=1}^{\mathsf{N}} \mathcal{R}_i^{\mathsf{T}} \mathcal{R}_i^{-1} \mathcal{R}_i,$$

where the linearized system is of the form

$$\mathcal{A} x := \begin{bmatrix} A_u & C_{uT} \\ C_{Tu} & A_T \end{bmatrix} \begin{bmatrix} x_u \\ x_T \end{bmatrix} = \begin{bmatrix} \tilde{r}_u \\ \tilde{r}_T \end{bmatrix} =: r.$$

For single-physics problems, we employ a **standard (R)GDSW preconditioner**.

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### Antarctica Velocity Problem – Reuse Strategies (Strong Scaling)

We employ different **reuse strategies** to **reduce the setup costs** of the two-level preconditioner

$$M_{\text{GDSW}}^{-1} = \boldsymbol{\Phi} \mathbf{K}_{\mathbf{0}}^{-1} \boldsymbol{\Phi}^{T} + \sum_{i=1}^{N} \mathbf{R}_{i}^{T} \mathbf{K}_{i}^{-1} \mathbf{R}_{i}.$$



	restriction operators			+ coarse basis			+ 0	oarse mat	rix
reuse	+ symboli	c fact. (1	st level)	+ symbol	lic fact.	(2nd level)			
MPI	avg. its	avg.	avg.	avg. its	avg	g. avg. its	avg. its	avg.	avg.
ranks	(nl its)	setup	solve	(nl its)	setu	p solve	(nl its)	setup	solve
512	<b>41.9</b> (11)	25.10 s	12.29 s	42.6 (11)	14.99	s 12.50 s	46.7 (11)	14.94 s	13.81 s
1 0 2 4	<b>43.3</b> (11)	9.18 s	5.85 s	44.5 (11)	5.65	<b>s</b> 6.08 s	49.2 (11)	5.75 s	6.78 s
2 0 4 8	<b>41.4</b> (11)	4.15 s	2.63 s	42.7 (11)	3.11	s 2.79 s	47.7 (11)	2.92 s	3.10 s
4 0 9 6	<b>41.2</b> (11)	1.66 s	1.49 s	42.5 (11)	1.07	s 1.54 s	48.9 (11)	0.95 s	1.75 s
8 1 9 2	<b>40.2</b> (11)	1.26 s	1.06 s	42.0 (11)	1.20	s 1.16 s	50.1 (11)	0.63 s	1.35 s
Problem	n: Velocity	Mesh:	Antarctica 4 km hor. 20 vert. la	a resolution ayers	Size:	35.3 m degrees of freedom (P1 FE)	s Coarse sp	ace: RGI	DSW

Cf. Heinlein, Perego, Rajamanickam (2022)

### **Greenland Coupled Problem – Coarse Spaces**

			fully coupled extensions					
			nc	reuse		reuse coarse basis		
	MPI		avg. its	avg.	avg.	avg. its	avg.	avg.
	ranks	dim $V_0$	(nl its)	setup	solve	(nl its)	setup	solve
	256	1 400	100.1 (27)	4.10 s	6.40 s	<b>18.5</b> (70)	2.28 s	1.07 s
	512	2852	129.1 (28)	1.88 s	4.20 s	<b>24.6</b> (38)	1.04 s	0.70 s
	1024	6 0 3 6	191.2 (65)	1.21 s	4.76 s	<b>34.2</b> (32)	0.66 s	0.70 s
	2048	12 368	237.4 (30)	0.96 s	4.06 s	<b>37.3</b> (30)	0.60 s	0.58 s
			deo	coupled	extension	IS		
			nc	reuse		reuse coarse basis		
	MPI		avg. its	avg.	avg.	avg. its	avg.	avg.
	ranks	dim $V_0$	(nl its)	setup	solve	(nl its)	setup	solve
	256	1 400	23.6 (29)	3.90 s	1.32 s	<b>21.5</b> (34)	2.23 s	1.18 s
	512	2852	27.5 (30)	1.83 s	0.78 s	<b>26.4</b> (33)	1.13 s	0.78 s
	1024	6 0 3 6	30.1 (29)	1.19 s	0.60 s	<b>28.6</b> (43)	0.66 s	0.61 s
	2048	12 368	36.4 (30)	0.69 s	0.56 s	<b>31.2</b> (50)	0.57 s	0.55 s
Problem:	Couple	d Mesh:	Greenland 3-30 km hor 20 vert. laye	. resolutio	Size:	7.5 m degr of freedom (P1 FE)	rees <b>Co</b> a	arse space:

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RGDSW

### **Greenland Coupled Problem – Large Problem**

|--|--|

	decoupled			fully coupled			decoupled		
	(n	o reuse)		(reuse coarse basis)			(reuse 1st level symb. fact.		
							+ co	barse basis)	)
MPI	avg. its.	avg.	avg.	avg. its	avg.	avg.	avg. its	avg.	avg.
ranks	(nl its)	setup	solve	(nl its)	setup	solve	(nl its)	setup	solve
512	<b>41.3</b> (36)	18.78 s	<b>4.99</b> s	45.3 (32)	11.84 s	5.35 s	45.0 (35)	<b>10.53</b> s	5.36 s
1024	53.0 (29)	8.68 s	4.22 s	<b>47.8</b> (37)	5.36 s	3.82 s	54.3 (32)	4.59 s	4.31 s
2 0 4 8	62.2 (86)	4.47 s	4.23 s	66.7 (38)	2.81 s	4.53 s	<b>59.1</b> (38)	2.32 s	3.99 s
4 0 9 6	<b>68.9</b> (40)	2.52 s	2.86 s	79.1 (36)	1.61 s	3.30 s	78.7 (38)	1.37 s	3.30 s
Problem	n: Coupled	Mesh:	Greenlar 1-10 km 20 vert.	nd hor. resolutio layers	Size:	68.6 m c of freedo (P1 FE)	legrees <b>Coa</b> om	rse space:	RGDSW

Cf. Heinlein, Perego, Rajamanickam (2022)

## Sparse Triangular Solver in KokkosKernels (Amesos2 – SuperLU/CHOLMOD)

The sparse triangular solver is an **important kernel** in many codes (including FROSch) but is **challenging to parallelize** 

- Factorization using a **sparse direct solver** typically leads to triangular matrices with **dense blocks** called **supernodes**
- In supernodal triangular solvers, rows/columns with a similar sparsity pattern are merged into a supernodal block, and the solve is then performed block-wise
- The parallelization potential for the triangular solver is determined by the sparsity pattern

Parallel supernode-based triangular solver:

- 1. Supernode-based level-set scheduling, where all leaf-supernodes within one level are solved in parallel (batched kernels for hierarchical parallelism)
- 2. Partitioned inverse of the submatrix associated with each level: SpTRSV is transformed into a sequence of SpMVs

See Yamazaki, Rajamanickam, and Ellingwood (2020) for more details.



Lower-triangular matrix – SuperLU

with METIS nested dissection ordering



### Three-Dimensional Linear Elasticity – Weak Scalability



Computations on Summit (OLCF): 42 IBM Power9 CPU cores and 6 NVIDIA V100 GPUs per node. Yamazaki, Heinlein, Rajamanickam (2023)

Alexander Heinlein (TU Delft)

### Three-Dimensional Linear Elasticity – Weak Scalability



# nodes	1	2	4	8	16			
# dofs	375K	750K	1.5M	3M	6M			
SUPERLU solve								
CPUs	2.03 (75)	2.07 (69)	1.87 (61)	1.95 (58)	2.48 (69)			
$n_p/\text{gpu} = 1$	1.43 (47)	1.52 (53)	2.82 (77)	2.44 (68)	2.61 (75)			
2	1.03 (46)	1.36 (65)	1.37 (60)	1.52 (65)	1.98 (86)			
4	0.93 (59)	0.91 (53)	0.98 (59)	1.33 (77)	1.21 (66)			
6	0.67 (46)	0.99 (65)	0.92 (57)	0.91 (57)	0.95 (57)			
7	1.03 (75)	1.04 (69)	0.90 (61)	0.97 (58)	1.18 (69)			
speedup	<b>2</b> .0×	<b>2</b> .0×	<b>2</b> .1×	<b>2</b> .0×	<b>2</b> .1×			
		Tacho	solve					
CPUs	1.60 (75)	1.63 (69)	1.49 (61)	1.51 (58)	1.90 (69)			
$n_p/\text{gpu} = 1$	1.17 (47)	1.37 (53)	1.92 (77)	1.78 (68)	2.21 (75)			
2	0.79 (46)	1.14 (65)	1.05 (60)	1.18 (65)	1.70 (86)			
4	0.85 (59)	0.81 (53)	0.78 (59)	1.22 (77)	1.19 (66)			
6	0.60 (46)	0.86 (65)	0.75 (57)	0.84 (57)	0.91 (57)			
7	0.99 (75)	0.93 (69)	0.82 (61)	0.93 (58)	1.22 (69)			
speedup	<b>1.6</b> ×	<b>1.8</b> ×	1.8  imes	1.6  imes	1.6 imes			

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### Three-Dimensional Linear Elasticity – ILU Subdomain Solver

ILU	J level	0	1	2	3					
	setup									
$\cap$	No	1.5	1.9	3.0	4.8					
G	ND	1.6	2.6	4.4	7.4					
	KK(No)	1.4	1.5	1.8	2.4					
	KK(ND)	1.7	2.0	2.9	5.2					
GF	Fast(No)	1.5	1.6	2.1	3.2					
	Fast(ND)	1.5	1.7	2.5	4.5					
spe	eedup	1.0×	1.2×	<b>1</b> .4×	1.5×					
			solve							
$\Box$	No	2.55 (158)	3.60 (112)	5.28 (99)	6.85 (88)					
ß	ND	4.17 (227)	5.36 (134)	6.61 (105)	7.68 (88)					
	KK(No)	3.81 (158)	4.12 (112)	4.77 (99)	5.65 (88)					
	KK(ND)	2.89 (227)	4.27 (134)	5.57 (105)	6.36 (88)					
5	Fast(No)	1.14 (173)	1.11 (141)	1.26 (134)	1.43 (126)					
	Fast(ND)	1.49 (227)	1.15 (137)	1.10 (109)	1.22 (100)					
spe	eedup	2.2×	3.2×	<b>4</b> .3×	<b>4</b> .8×					

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### **ILU** variants

- Kokkoskernels ILU (KK)
- FASTILU (Fast); cf. Chow, Patel (2015) and Boman, Patel, Chow, Rajamanickam (2016)

No reordering (No) and nested dissection (ND)  $% \left( ND\right) =0$ 



### Three-Dimensional Linear Elasticity – Weak Scalability Using ILU

# r	nodes	1	2	4	8	16					
# c	lofs	648 K	1.2 M	2.6 M	5.2 M	10.3 M					
	setup										
CP	U	1.9	2.2	2.4	2.4	2.6					
Ú	KK	1.4	2.0	2.2	2.4	2.8					
GP	Fast	1.5	2.2	2.3	2.5	2.8					
spe	edup	1.3×	<b>1.0</b> ×	<b>1.0</b> ×	<b>1.0</b> ×	<b>0.9</b> ×					
			SO	lve							
CP	U	3.60 (112)	7.26 (84)	6.93 (78)	6.41 (75)	4.1 (109)					
D	KK	4.3 (119)	3.9 (110)	4.8 (105)	4.3 (97)	4.9 (109)					
GF	Fast	1.2 (154)	1.0 (133)	1.1 (130)	1.3 (117)	1.6 (131)					
spe	edup	3.3×	<b>3.8</b> ×	<b>3.4</b> ×	<b>2.5</b> ×	<b>2.6</b> ×					

Computations on Summit (OLCF): 42 IBM Power9 CPU cores and 6 NVIDIA V100 GPUs per node.

Yamazaki, Heinlein, Rajamanickam (2023)

### Thank you for your attention!

### Summary

- FROSCH is based on the Schwarz framework and energy-minimizing coarse spaces, which provide numerical scalability using only algebraic information for a variety of applications including nonlinear multi-physics problems
- For nonlinear problems,
  - the reuse of components of the preconditioner and
  - the speedup of the solver phase (e.g., using GPUs)

can significantly help to improve the solver performance.

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# Thank you for your attention!

- $\rightarrow$  Talk by Ichitaro Yamazaki on Tuesday (00911 (2/2)): Related nonlinear Schwarz methods
- $\rightarrow$  Talk by Martin Lanser on Thursday (01054 (2/3)): FROSch on GPUs
- $\rightarrow$  Talk by Friederike Röver on Friday (01054 (3/3)): FROSch for chemo-mechanics problems