



Fast and Robust Overlapping Schwarz (FROSch) Domain Decomposition Preconditioners

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Solving A Model Problem



Consider a diffusion model problem:

 $-\nabla \cdot (\alpha(x)\nabla u(x)) = f \quad \text{in } \Omega = [0, 1]^2,$ $u = 0 \quad \text{on } \partial\Omega.$

Discretization using finite elements yields a **sparse** linear system of equations

Ku = f.

Direct solvers

For fine meshes, solving the system using a direct solver is not feasible due to **superlinear complexity and memory cost**.

Iterative solvers

Iterative solvers are efficient for solving sparse linear systems of equations, however, the convergence rate generally depends on the condition number κ (*A*). It deteriorates, e.g., for

- fine meshes, that is, small element sizes h
- large contrasts $\frac{\max_{x} \alpha(x)}{\min_{x} \alpha(x)}$

 \Rightarrow We introduce a preconditioner $M^{-1} \approx A^{-1}$ to improve the condition number:

$$\boldsymbol{M}^{-1}\boldsymbol{A}\boldsymbol{u}=\boldsymbol{M}^{-1}\boldsymbol{f}$$

Two-Level Schwarz Preconditioners

One-level Schwarz preconditioner





Based on an overlapping domain decomposition, we define a one-level Schwarz operator

$$\boldsymbol{M}_{\text{OS-1}}^{-1}\boldsymbol{K} = \sum_{i=1}^{N} \boldsymbol{R}_{i}^{\top}\boldsymbol{K}_{i}^{-1}\boldsymbol{R}_{i}\boldsymbol{K}_{i}$$

where \mathbf{R}_i and \mathbf{R}_i^{\top} are restriction and prolongation operators corresponding to Ω'_i , and $\mathbf{K}_i := \mathbf{R}_i \mathbf{K} \mathbf{R}_i^{\top}$.

Condition number estimate:

$$\kappa\left(\pmb{M}_{\mathsf{OS-1}}^{-1}\pmb{K}
ight) \leq C\left(1+rac{1}{H\delta}
ight)$$

with subdomain size H and overlap width δ .

Lagrangian coarse space





The two-level overlapping Schwarz operator reads

$$\boldsymbol{M}_{\text{OS-2}}^{-1}\boldsymbol{K} = \underbrace{\boldsymbol{\Phi}\boldsymbol{K}_{0}^{-1}\boldsymbol{\Phi}^{\top}\boldsymbol{K}}_{\text{coarse level - global}} + \underbrace{\sum_{i=1}^{N}\boldsymbol{R}_{i}^{\top}\boldsymbol{K}_{i}^{-1}\boldsymbol{R}_{i}\boldsymbol{K}}_{\text{first level - local}},$$

where Φ contains the coarse basis functions and $K_0 := \Phi^\top K \Phi$; cf., e.g., **Toselli, Widlund (2005)**. The construction of a Lagrangian coarse basis requires a coarse triangulation.

Condition number estimate:

$$\kappa\left(\boldsymbol{M}_{\mathsf{OS-2}}^{-1}\boldsymbol{K}
ight)\leq C\left(1+rac{\boldsymbol{H}}{\delta}
ight)$$

Two-Level Schwarz Preconditioners



FROSch (Fast and Robust Overlapping Schwarz) Framework in Trilinos





Software

- Object-oriented C++ domain decomposition solver framework with $\rm MPI\text{-}based$ distributed memory parallelization
- Part of TRILINOS with support for both parallel linear algebra packages EPETRA and TPETRA
- Node-level parallelization and performance portability on CPU and GPU architectures through KOKKOS and KOKKOSKERNELS
- Accessible through unified ${\rm TRILINOS}$ solver interface ${\rm STRATIMIKOS}$

Methodology

- Parallel scalable multi-level Schwarz domain decomposition
 preconditioners
- Algebraic construction based on the parallel distributed system matrix
- Extension-based coarse spaces

Team (active)

- Filipe Cumaru (TU Delft)
- Kyrill Ho (UCologne)
- Jascha Knepper (UCologne)
- Friederike Röver (TUBAF)
- Lea Saßmannshausen (UCologne)

- Alexander Heinlein (TU Delft)
- Axel Klawonn (UCologne)
- Siva Rajamanickam (SNL)
- Oliver Rheinbach (TUBAF)
- Ichitaro Yamazaki (SNL)

Overlapping domain decomposition

Adjacency can be determined algebraically from the sparsity pattern of the system matrix *A* (nonzero off-diagonal entries).



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Coarse space

1. Interface components





For scalar elliptic problems, the null space consists only of constant functions.





Overlapping domain decomposition Overlap $\delta = 1h$ **Overlap** $\delta = 2h$ Nonoverlapping DD Adjacency can be determined algebraically from the sparsity pattern of the system matrix A (nonzero off-diagonal entries). **Coarse space** 3. Extension 1. Interface components 2. Interface basis (partition of unity \times null space)

Examples of FROSch Coarse Spaces

GDSW (Generalized Dryja-Smith-Widlund)





- Dohrmann, Klawonn, Widlund (2008)
- Dohrmann, Widlund (2009, 2010, 2012)

MsFEM (Multiscale Finite Element Method)





- Hou (1997), Efendiev and Hou (2009)
- Buck, Iliev, and Andrä (2013)
- H., Klawonn, Knepper, Rheinbach (2018)

RGDSW (Reduced dimension GDSW)





- Dohrmann, Widlund (2017)
- H., Klawonn, Knepper, Rheinbach, Widlund (2022)

Q1 Lagrangian / piecewise bilinear





Piecewise linear interface partition of unity functions and a **structured domain decomposition**.

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Examples of FROSch Coarse Spaces



For elliptic model problems, the condition number of the (R)GDSW two-level Schwarz operator is bounded by

$$\kappa\left(\pmb{M}_{(\mathsf{R})\mathsf{GDSW}}^{-1}\pmb{K}\right) \leq C\left(1+\frac{H}{\delta}\right)\left(1+\log\left(\frac{H}{h}\right)\right)^{\alpha},$$

where

C constant (does not depend on h, H, or δ),

H subdomain diameter,

h element size,

 δ width of the overlap,

 $\alpha \in \{0, 1, 2\}$ power (depends on problem dimension, regularity of the subdomains, and variant of the algorithm).

Weak Scalability up to 64 k MPI Ranks / 1.7 b Unknowns (3D Poisson; Juqueen)

GDSW vs RGDSW (reduced dimension)





Two-level vs three-level GDSW

Heinlein, Klawonn, Rheinbach, Röver (2019, 2020).



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1000

Cores

10000

100000

70

35

0

100

Monolithic and Adaptive Extension-Based Coarse Spaces

Monolithic (R)GDSW Preconditioners for CFD Simulations

Consider the discrete saddle point problem

$$\mathcal{A}_{X} = \begin{bmatrix} \mathbf{K} & \mathbf{B}^{\top} \\ \mathbf{B} & \mathbf{0} \end{bmatrix} \begin{bmatrix} \mathbf{u} \\ \mathbf{p} \end{bmatrix} = \begin{bmatrix} \mathbf{f} \\ \mathbf{0} \end{bmatrix} = \mathbf{6}.$$

Monolithic GDSW preconditioner

We construct a monolithic GDSW preconditioner

$$\mathcal{M}_{\mathsf{GDSW}}^{-1} = \phi \mathcal{R}_0^{-1} \phi^\top + \sum\nolimits_{i=1}^N \mathcal{R}_i^\top \mathcal{R}_i^{-1} \mathcal{R}_i$$

with block matrices $\mathcal{A}_0=\phi^\top\mathcal{A}\phi,\,\mathcal{A}_i=\mathcal{R}_i\mathcal{A}\mathcal{R}_i^\top,$ and

$$\mathcal{R}_i = \begin{bmatrix} \mathcal{R}_{u,i} & \mathbf{0} \\ \mathbf{0} & \mathcal{R}_{p,i} \end{bmatrix} \quad \text{and} \quad \phi = \begin{bmatrix} \Phi_{u,u_0} & \Phi_{u,p_0} \\ \Phi_{p,u_0} & \Phi_{p,p_0} \end{bmatrix}.$$

Using \mathcal{A} to compute extensions: $\phi_I = -\mathcal{A}_{II}^{-1}\mathcal{A}_{I\Gamma}\phi_{\Gamma}$; cf. Heinlein, Hochmuth, Klawonn (2019, 2020).







Stokes flow

Navier-Stokes flow

Related work:

- Original work on monolithic Schwarz preconditioners: Klawonn and Pavarino (1998, 2000)
- Other publications on monolithic Schwarz preconditioners: e.g., Hwang and Cai (2006), Barker and Cai (2010), Wu and Cai (2014), and the presentation Dohrmann (2010) at the Workshop on Adaptive Finite Elements and Domain Decomposition Methods in Milan.

Monolithic (R)GDSW Preconditioners for CFD Simulations

Consider the discrete saddle point problem

$$\mathcal{A}_{X} = \begin{bmatrix} \mathbf{K} & \mathbf{B}^{\top} \\ \mathbf{B} & \mathbf{0} \end{bmatrix} \begin{bmatrix} \mathbf{u} \\ \mathbf{p} \end{bmatrix} = \begin{bmatrix} \mathbf{f} \\ \mathbf{0} \end{bmatrix} = \mathbf{6}.$$

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We construct a monolithic GDSW preconditioner

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with block matrices $\mathcal{A}_0 = \phi^\top \mathcal{A} \phi$, $\mathcal{A}_i = \mathcal{R}_i \mathcal{A} \mathcal{R}_i^\top$.

SIMPLE block preconditioner

We employ the SIMPLE (Semi-Implicit Method for Pressure Linked Equations) block preconditioner

$$\mathcal{M}_{\mathsf{SIMPLE}}^{-1} = \begin{bmatrix} \mathbf{I} & -\mathbf{D}^{-1}\mathbf{B} \\ \mathbf{0} & \alpha \mathbf{I} \end{bmatrix} \begin{bmatrix} \mathbf{K}^{-1} & \mathbf{0} \\ -\hat{\mathbf{S}}^{-1}\mathbf{B}\mathbf{K}^{-1} & \hat{\mathbf{S}}^{-1} \end{bmatrix};$$

see Patankar and Spalding (1972). Here,

- $\hat{\boldsymbol{S}} = -\boldsymbol{B}\boldsymbol{D}^{-1}\boldsymbol{B}^{\top}$, with $\boldsymbol{D} = \operatorname{diag} \boldsymbol{K}$
- α is an under-relaxation parameter

We **approximate the inverses** using (R)GDSW preconditioners.

Monolithic vs. SIMPLE preconditioner



Steady-state Navier-Stokes equations

prec.	# MPI ranks	243	1 1 2 5	15 562
Monolithic	setup	39.6 s	57.9 s	95.5 s
RGDSW	solve	57.6 s	69.2 s	74.9 s
(FROSCH)	total	97.2 s	127.7 s	170.4 s
(FROSCH) SIMPLE	total setup	97.2 s 39.2 s	127.7 s 38.2 s	170.4 s 68.6 s
(FROSCH) SIMPLE RGDSW (Теко	total setup solve	97.2 s 39.2 s 86.2 s	127.7 s 38.2 s 106.6 s	170.4 s 68.6 s 127.4 s

Computations on Piz Daint (CSCS). Implementation in the finite element software FEDDLib.

Balancing the Velocity and Pressure Coarse Spaces



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Local Pressure Projections

We slightly modify the monolithic two-level overlapping Schwarz preconditioner

$$\mathcal{M}_{\mathsf{OS-2}}^{-1} = \phi \mathcal{A}_0^{-1} \phi^\top + \sum_{i=1}^N \mathcal{R}_i^\top \overline{\mathcal{P}}_i \mathcal{A}_i^{-1} \mathcal{R}_i$$

with local projection operators $\overline{\mathcal{P}}_i$ of the form

$$\overline{\mathcal{P}}_{i} = \begin{bmatrix} I_{u,i} & 0\\ 0 & \overline{P}_{p,i} \end{bmatrix}, \text{ with } \overline{P}_{p,i} = I_{p,i} - \frac{a_{i}a_{i}^{T}}{a_{i}^{T}a_{i}},$$

where a_i is the discretization of the integral $\int_{\Omega_i} u \, dx$.

Backward-facing step (Re = 200)



Lid-driven cavity Stokes



Navier–Stokes flow with kinematic viscosity ν =1e^{-3}

♯ MPI ran	iks	196	1089	4 356
P2-P1,	Local projections	33	36	35
H/h = 50	—	118	172	237
$\begin{array}{l} P1-P1stab \\ H/h = 80 \end{array}$	Pressure stab.	38	36	36

RGDSW-RGDSW coarse spaces

 \rightarrow Local pressure corrections significantly improve the convergence. We obtain very fast convergence using a good combination of coarse spaces.

Heinlein, Klawonn, Saßmannshausen (in prep.)

Computations on Fritz (FAU).

FROSch Preconditioners for Land Ice Simulations



https://github.com/SNLComputation/Albany

The velocity of the ice sheet in Antarctica and Greenland is modeled by a **first-order-accurate Stokes approximation model**,

$$-\nabla \cdot (2\mu \dot{\epsilon}_1) + \rho g \frac{\partial s}{\partial x} = 0, \quad -\nabla \cdot (2\mu \dot{\epsilon}_2) + \rho g \frac{\partial s}{\partial y} = 0.$$



with a nonlinear viscosity model (Glen's law); cf., e.g., Blatter (1995) and Pattyn (2003).

	Antarctica (velocity)			Greenland	(multiphysics v	el. & temperature)
	4 km resolution, 20 layers, 35 m dofs			1-10 km resolution, 20 layers, 69 m dofs		
MPI ranks	avg. its	avg. setup	avg. solve	avg. its	avg. setup	avg. solve
512	41.9 (11)	25.10 s	12.29 s	41.3 (36)	18.78 s	4.99 s
1 024	43.3 (11)	9.18 s	5.85 s	53.0 (29)	8.68 s	4.22 s
2 048	41.4 (11)	4.15 s	2.63 s	62.2 (86)	4.47 s	4.23 s
4 096	41.2 (11)	1.66 s	1.49 s	68.9 (40)	2.52 s	2.86 s
8 192	40.2 (11)	1.26 s	1.06 s	-	-	-

Computations performed on Cori (NERSC).

Heinlein, Perego, Rajamanickam (2022)

Adaptive Extension-Based Coarse Spaces for Schwarz Preconditioners

Highly heterogeneous problems ...

... appear in most areas of modern science and engineering:







Micro section of a dual-phase steel. Courtesy of J. Schröder.

Groundwater flow (SPE10); cf. Christie and Blunt (2001). Composition of arterial walls; taken from **O'Connell et al. (2008)**.

Adaptive coarse spaces

The coarse space is **enhanced** by eigenfunctions of **local edge and face eigenvalue problems** with eigenvalues below tolerances $tol_{\mathcal{E}}$ and $tol_{\mathcal{F}}$:

$$\kappa\left(\mathbf{M}_{*}^{-1}\mathbf{K}\right) \leq C\left(1 + \frac{1}{\operatorname{tol}_{\mathcal{B}}} + \frac{1}{\operatorname{tol}_{\mathcal{F}}} + \frac{1}{\operatorname{tol}_{\mathcal{F}} \cdot \operatorname{tol}_{\mathcal{F}}}\right);$$

C does not depend on *h*, *H*, or the coefficients. OS-ACMS & adaptive GDSW (AGDSW) (Heinlein, Klawonn, Knepper, Rheinbach (2018, 2018, 2019)).

Local eigenvalue problems

Local generalized eigenvalue problems corresponding to the edges ${\mathcal S}$ and faces ${\mathcal F}$ of the domain decomposition:

$$\begin{aligned} \forall E \in \mathcal{E} : \qquad & \boldsymbol{S}_{EE} \boldsymbol{\tau}_{*,E} = \lambda_{*,E} \boldsymbol{K}_{EE} \boldsymbol{\tau}_{*,E}, \quad \forall \boldsymbol{\tau}_{*,E} \in \boldsymbol{V}_{E}, \\ \forall F \in \mathcal{F} : \qquad & \boldsymbol{S}_{FE} \boldsymbol{\tau}_{*,F} = \lambda_{*,E} \boldsymbol{K}_{FE} \boldsymbol{\tau}_{*,F}, \quad \forall \boldsymbol{\tau}_{*,F} \in \boldsymbol{V}_{F}, \end{aligned}$$

with Schur complements S_{EE} , S_{FF} with Neumann boundary conditions and submatrices K_{EE} , K_{FF} of K. We select eigenfunctions corresponding to eigenvalues below tolerances $tol_{\&}$ and $tol_{\mathcal{J}}$.

 \rightarrow The corresponding coarse basis functions are **energy-minimizing extensions** into the interior of the subdomains.



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C does not depend on *h*, *H*, or the coefficients. OS-ACMS & adaptive GDSW (AGDSW) (Heinlein, Klawonn, Knepper, Rheinbach (2018, 2018, 2019)).

Foam coefficient function example



Solid phase: $A = 10^6$; transparent phase: A = 1; 100 subdomains

V ₀	tol ₈	$\mathit{tol}_{\mathcal{F}}$	it.	κ	dim V_0	$\dim V_0/\operatorname{dof}$
$V_{ m GDSW}$	—	_	565	$1.3 \cdot 10^{6}$	1601	0.27 %
$V_{ m AGDSW}$	0.05	0.05	60	30.2	1968	0.33 %
$V_{\rm OS-ACMS}$	0.001	0.001	57	30.3	690	0.12 %

Cf. Heinlein, Klawonn, Knepper, Rheinbach (2018, 2019).

Algebraic Adaptive Extension-Based Coarse Spaces

Two algebraic eigenvalue problems

Use the a-orthogonal decomposition

 $V_{\Omega_{e}} = V_{\Omega_{e}}^{0} \oplus \{E_{\partial \Omega_{e} \to \Omega_{e}}(v) : v \in V_{\partial \Omega_{e}}\}$

to "split the AGDSW (Neumann) eigenvalue problem" into two:

- Dirichlet eigenvalue problem on V⁰_{Ω_e}
- Transfer eigenvalue problem on $V_{\Omega_{e},harm}$; cf. Smetana, Patera (2016)



Condition number estimate

$$\kappa\left(\boldsymbol{M}_{\mathsf{DIR}\&\mathsf{TR}}^{-1}\boldsymbol{K}
ight)\leq C\max\left\{1/ au \mathsf{OL}_{\mathsf{DIR}}, \ au \mathsf{OL}_{\mathsf{TR}}/lpha_{\mathsf{min}}
ight\},$$

where *C* is independent of *H*, *h*, and the contrast of the coefficient function α .

Heinlein & Smetana (subm. 2023; preprint arXiv).

Numerical results – SPE10 benchmark

Layer 70 from model 2; cf. Christie and Blunt (2001)



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Accelerating Time-to-Solution

Inexact Subdomain Solvers in FROSch

$$oldsymbol{M}_{ ext{OS-2}}^{-1}oldsymbol{\mathcal{K}}=\Phioldsymbol{\mathcal{K}}_0^{-1}\Phi^{ op}oldsymbol{\mathcal{K}}+\sum_{i=1}^Noldsymbol{R}_i^{ op}oldsymbol{\mathcal{K}}_i^{-1}oldsymbol{R}_ioldsymbol{\mathcal{K}}$$

3D Laplacian; 512 MPI ranks = 512 (= 8 \times 8 \times 8) subdomains; H/δ = 10; RGDSW coarse space.

		subdomain solver							
		direct	ILU	l(k)	symm. G	auß–Seidel	Chebyshev polyn.		
		solver	k = 2	k = 3	5 sweeps	10 sweeps	p = 6	p = 8	
H/h = 20	iter	26	33	30	31	28	34	31	
H/H = 20,	setup time	1.89 s	0.97 s	1.01 s	0.89 s	0.91 s	0.73 s	0.71 s	
\approx 14 K dots	apply time	0.39 s	0.27 s	0.31 s	0.31 s	0.35 s	0.30 s	0.30 s	
per rank	prec. time	2.28 s	1.24 s	1.32 s	1.20 s	1.26 s	1.03 s	1.01 s	
H/h = 40	iter	30	55	46	52	41	59	51	
H/H = 40,	setup time	12.09 s	6.14 s	6.26 s	5.74 s	5.89 s	5.55 s	5.64 s	
$\approx 105 \text{k}$ dois	apply time	4.21 s	1.84 s	1.96 s	2.66 s	3.28 s	2.52 s	2.47 s	
per rank	prec. time	16.30 s	7.98 s	8.22 s	8.40 s	9.18 s	8.16 s	8.11 s	
H/h = 60	iter		81	64	76	56	88	74	
n/n = 00,	setup time	0.0M	47.29 s	47.87 s	45.14 s	45.08 s	45.44 s	45.49 s	
	apply time	00101	10.79 s	9.98 s	13.00 s	16.16 s	11.95 s	12.09 s	
регланк	prec. time		58.08 s	57.85 s	58.15 s	61.25 s	57.39 s	57.59 s	

INTEL MKL PARDISO; ILU / symmetric Gauß–Seidel / Chebyshev polynomials from IFPACK2.

Parallel computations on dual-socket Intel Xeon Platinum machine at Sandia National Laboratories (Blake).

Inexact Subdomain Solvers in FROSch

$$oldsymbol{M}_{ ext{OS-2}}^{-1}oldsymbol{\mathcal{K}}=\Phioldsymbol{\mathcal{K}}_0^{-1}\Phi^{ au}oldsymbol{\mathcal{K}}+\sum_{i=1}^Noldsymbol{R}_i^{ au}oldsymbol{\mathsf{K}}_i^{-1}oldsymbol{R}_ioldsymbol{\mathcal{K}}$$

3D Laplacian; 512 MPI ranks = 512 (= 8 \times 8 \times 8) subdomains; H/δ = 10; RGDSW coarse space.

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$\approx 14 \text{K}$ dois	apply time	0.39 s	0.27 s	0.31 s	0.31 s	0.35 s	0.30 s	0.30 s
per rank	prec. time	2.28 s	1.24 s	1.32 s	1.20 s	1.26 s	1.03 s	1.01 s
H/h = 40	iter	30	55	46	52	41	59	51
H/H = 40,	setup time	12.09 s	6.14 s	6.26 s	5.74 s	5.89 s	5.55 s	5.64 s
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H/h = 60	iter		81	64	76	56	88	74
11/11 = 00,	setup time	0.014	47.29 s	47.87 s	45.14 s	45.08 s	45.44 s	45.49 s
	apply time	00101	10.79 s	9.98 s	13.00 s	16.16 s	11.95 s	12.09 s
регланк	prec. time		58.08 s	57.85 s	58.15 s	61.25 s	57.39 s	57.59 s

INTEL MKL PARDISO; ILU / symmetric Gauß–Seidel / Chebyshev polynomials from IFPACK2.

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Inexact Extension Solvers in FROSch

$$\Phi = \begin{bmatrix} -\mathbf{K}_{II}^{-1}\mathbf{K}_{\Gamma I}^{T}\Phi_{\Gamma} \\ \Phi_{\Gamma} \end{bmatrix} = \begin{bmatrix} \Phi_{I} \\ \Phi_{\Gamma} \end{bmatrix}$$

3D Laplacian; 512 MPI ranks = 512 (= 8 \times 8 \times 8) subdomains; H/δ = 10; RGDSW coarse space.

extension	solver	direct		precond	itioned GMF	RES (rel. tol.	$= 10^{-4}$)	
(10 Gauss–Seide	el sweeps for	antect	ILU(k)		symm. G	auß–Seidel	Chebyshev polyn.	
the subdomain solver)		solver	k = 2	k = 3	5 sweeps	10 sweeps	p = 6	p = 8
H/h = 20	iter	28	28	28	28	28	28	28
H/H = 20,	setup time	0.89 s	0.93 s	0.89 s	0.78 s	0.83 s	0.79 s	0.84 s
$\approx 14 \text{K}$ dois	apply time	0.35 s	0.35 s	0.34 s	0.36 s	0.34 s	0.35 s	0.34 s
per rank	prec. time	1.23 s	1.28 s	1.23 s	1.14 s	1.17 s	1.14 s	1.18 s
H/h = 40	iter	41	41	41	41	41	41	41
H/H = 40,	setup time	5.72 s	4.16 s	4.61 s	4.26 s	4.64 s	4.27 s	4.33 s
$\approx 105 \text{k}$ dois	apply time	3.33 s	3.33 s	3.30 s	3.33 s	3.30 s	3.28 s	3.29 s
per rank	prec. time	9.04 s	7.49 s	7.92 s	7.59 s	7.95 s	7.55 s	7.62 s
H/h = 60	iter	56	56	56	56	56	56	56
n/n = 00,	setup time	45.16 s	17.75 s	18.16 s	17.98 s	19.34 s	17.93 s	18.04 s
	apply time	15.83 s	18.04 s	17.08 s	16.26 s	15.81 s	16.19 s	16.44 s
per rank	prec. time	60.99 s	35.79 s	35.25 s	34.24 s	35.15 s	34.12 s	34.49 s

INTEL MKL PARDISO; ILU / symmetric Gauß–Seidel / Chebyshev polynomials from $\rm IFPACK2.$

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$$\Phi = \begin{bmatrix} -\mathbf{K}_{\mathsf{II}}^{-1}\mathbf{K}_{\mathsf{\Gamma}^{\prime}}^{\mathsf{T}}\Phi_{\mathsf{\Gamma}} \\ \Phi_{\mathsf{\Gamma}} \end{bmatrix} = \begin{bmatrix} \Phi_{\mathsf{I}} \\ \Phi_{\mathsf{\Gamma}} \end{bmatrix}$$

3D Laplacian; 512 MPI ranks = 512 (= 8 \times 8 \times 8) subdomains; H/δ = 10; RGDSW coarse space.

extension	solver	direct		precond	itioned GMF	RES (rel. tol.	$= 10^{-4}$)	
(10 Gauss–Seide	el sweeps for	antect	ILU(k)		symm. Gauß–Seidel		Chebyshev polyn.	
the subdomain solver)		solver	k = 2	k=3	5 sweeps	10 sweeps	p = 6	p = 8
11/1- 20	iter	28	28	28	28	28	28	28
H/H = 20,	setup time	0.89 s	0.93 s	0.89 s	0.78 s	0.83 s	0.79 s	0.84 s
\approx 14 K dots	apply time	0.35 s	0.35 s	0.34 s	0.36 s	0.34 s	0.35 s	0.34 s
per rank	prec. time	1.23 s	1.28 s	1.23 s	1.14 s	1.17 s	1.14 s	1.18 s
H/h = 40	iter	41	41	41	41	41	41	41
11/11 = 40, $\approx 105 k dots$	setup time	5.72 s	4.16 s	4.61 s	4.26 s	4.64 s	4.27 s	4.33 s
$\approx 105 \text{ k uois}$	apply time	3.33 s	3.33 s	3.30 s	3.33 s	3.30 s	3.28 s	3.29 s
per rank	prec. time	9.04 s	7.49 s	7.92 s	7.59 s	7.95 s	7.55 s	7.62 s
H/h = 60	iter	56	56	56	56	56	56	56
n/n = 00,	setup time	45.16 s	17.75 s	18.16 s	17.98 s	19.34 s	17.93 s	18.04 s
\sim 330 k dois	apply time	15.83 s	18.04 s	17.08 s	16.26 s	15.81 s	16.19 s	16.44 s
per rank	prec. time	60.99 s	35.79 s	35.25 s	34.24 s	35.15 s	34.12 s	34.49 s

INTEL MKL PARDISO; ILU / symmetric Gauß–Seidel / Chebyshev polynomials from $\rm IFPACK2.$

Parallel computations on dual-socket Intel Xeon Platinum machine at Sandia National Laboratories (Blake).

Sparse Triangular Solver in KokkosKernels (Amesos2 – SuperLU/Tacho)

per node.

SuperLU & SpTRSV

- Supernodal LU factorization with partial pivoting
- Triangular solver with level-set scheduling (KOKKOSKERNELS);
 - cf. Yamazaki, Rajamanickam,

Ellingwood (2020).



Tacho

- Multifrontal factorization with pivoting inside frontal matrices
- Implementation using KOKKOS using level-set scheduling
- Cf. Kim, Edwards, Rajamanickam (2018).



Three-Dimensional Linear Elasticity – Weak Scalability



# nodes	1	2	4	8	16			
# dofs	375 K	750 K	1.5 M	3 M	6 M			
	SuperLU solve							
CPUs	2.03 (75)	2.07 (69)	1.87 (61)	1.95 (58)	2.48 (69)			
$n_p/\text{GPU} = 1$	1.43 (47)	1.52 (53)	2.82 (77)	2.44 (68)	2.61 (75)			
2	1.03 (46)	1.36 (65)	1.37 (60)	1.52 (65)	1.98 (86)			
4	0.93 (59)	0.91 (53)	0.98 (59)	1.33 (77)	1.21 (66)			
6	0.67 (46)	0.99 (65)	0.92 (57)	0.91 (57)	0.95 (57)			
7	1.03 (75)	1.04 (69)	0.90 (61)	0.97 (58)	1.18 (69)			
speedup	2.0×	2 .0×	2 .1×	2 .0×	2 .1×			
speedup	2.0 ×	2.0×	2.1×	2.0 ×	2 .1×			
speedup CPUs	2.0×	2.0× Tacho 1.63 (69)	2.1× solve 1.49 (61)	2.0× 1.51 (58)	2.1× 1.90 (69)			
speedup CPUs n _p /GPU = 1	2.0× 1.60 (75) 1.17 (47)	2.0× TACHO 1.63 (69) 1.37 (53)	2.1× solve 1.49 (61) 1.92 (77)	2.0× 1.51 (58) 1.78 (68)	2.1× 1.90 (69) 2.21 (75)			
speedup CPUs $n_p/GPU = 1$ 2	2.0× 1.60 (75) 1.17 (47) 0.79 (46)	2.0× TACHO 1.63 (69) 1.37 (53) 1.14 (65)	2.1× solve 1.49 (61) 1.92 (77) 1.05 (60)	2.0× 1.51 (58) 1.78 (68) 1.18 (65)	2.1× 1.90 (69) 2.21 (75) 1.70 (86)			
speedup CPUs $n_p/GPU = 1$ 2 4	2.0× 1.60 (75) 1.17 (47) 0.79 (46) 0.85 (59)	2.0× TACHO 1.63 (69) 1.37 (53) 1.14 (65) 0.81 (53)	2.1× solve 1.49 (61) 1.92 (77) 1.05 (60) 0.78 (59)	2.0× 1.51 (58) 1.78 (68) 1.18 (65) 1.22 (77)	2.1× 1.90 (69) 2.21 (75) 1.70 (86) 1.19 (66)			
speedup CPUs $n_p/GPU = 1$ 2 4 6	2.0× 1.60 (75) 1.17 (47) 0.79 (46) 0.85 (59) 0.60 (46)	2.0× TACHO 1.63 (69) 1.37 (53) 1.14 (65) 0.81 (53) 0.86 (65)	2.1× solve 1.49 (61) 1.92 (77) 1.05 (60) 0.78 (59) 0.75 (57)	2.0× 1.51 (58) 1.78 (68) 1.18 (65) 1.22 (77) 0.84 (57)	2.1× 1.90 (69) 2.21 (75) 1.70 (86) 1.19 (66) 0.91 (57)			
speedup CPUs $n_p/GPU = 1$ 2 4 6 7	2.0× 1.60 (75) 1.17 (47) 0.79 (46) 0.85 (59) 0.60 (46) 0.99 (75)	2.0× TACHO 1.63 (69) 1.37 (53) 1.14 (65) 0.81 (53) 0.86 (65) 0.93 (69)	2.1× solve 1.49 (61) 1.92 (77) 1.05 (60) 0.78 (59) 0.75 (57) 0.82 (61)	2.0× 1.51 (58) 1.78 (68) 1.18 (65) 1.22 (77) 0.84 (57) 0.93 (58)	2.1× 1.90 (69) 2.21 (75) 1.70 (86) 1.19 (66) 0.91 (57) 1.22 (69)			

Computations on Summit (OLCF): 42 IBM Power9 CPU cores and 6 NVIDIA V100 GPUs per node.

Yamazaki, Heinlein, Rajamanickam (2023)

Alexander Heinlein (Delft University of Technology)

Three-Dimensional Linear Elasticity – ILU Subdomain Solver

ILU	J level	0	1	2	3		
	setup						
Ū	No	1.5	1.9	3.0	4.8		
G	ND	1.6	2.6	4.4	7.4		
	KK(No)	1.4	1.5	1.8	2.4		
	KK(ND)	1.7	2.0	2.9	5.2		
В	Fast(No)	1.5	1.6	2.1	3.2		
	Fast(ND)	1.5	1.7	2.5	4.5		
spe	eedup	1.0×	1 .2×	1 .4×	1.5 imes		
			solve				
∩.	No	2.55 (158)	3.60 (112)	5.28 (99)	6.85 (88)		
Ы	ND	4.17 (227)	5.36 (134)	6.61 (105)	7.68 (88)		
	KK(No)	3.81 (158)	4.12 (112)	4.77 (99)	5.65 (88)		
	KK(ND)	2.89 (227)	4.27 (134)	5.57 (105)	6.36 (88)		
5	Fast(No)	1.14 (173)	1.11 (141)	1.26 (134)	1.43 (126)		
	Fast(ND)	1.49 (227)	1.15 (137)	1.10 (109)	1.22 (100)		
spe	eedup	2.2×	3.2×	4 .3×	4.8 ×		

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Yamazaki, Heinlein, Rajamanickam (2023)

ILU variants

- KokkosKernels ILU (KK)
- Iterative FASTILU (Fast); cf. Chow, Patel (2015) and Boman, Patel, Chow, Rajamanickam (2016)

No reordering (\mbox{No}) and nested dissection (\mbox{ND})





total time

Three-Dimensional Linear Elasticity – Weak Scalability Using ILU(1)

# r	odes	1	2	4	8	16		
# c	lofs	648 K	1.2 M	2.6 M	5.2 M	10.3 M		
	setup							
CP	U	1.9	2.2	2.4	2.4	2.6		
Ŋ	KK	1.4	2.0	2.2	2.4	2.8		
GF	Fast	1.5	2.2	2.3	2.5	2.8		
spe	edup	1.3×	1.0 imes	1 .0×	1 .0×	0.9 ×		
			sol	ve				
CP	U	3.60 (112)	7.26 (84)	6.93 (78)	6.41 (75)	4.1 (109)		
Ŋ	KK	4.3 (119)	3.9 (110)	4.8 (105)	4.3 (97)	4.9 (109)		
GF	Fast	1.2 (154)	1.0 (133)	1.1 (130)	1.3 (117)	1.6 (131)		
spe	edup	3 .3×	3.8 ×	3 .4×	2 .5×	2.6 ×		

Computations on Summit (OLCF): 42 IBM Power9 CPU cores and 6 NVIDIA V100 GPUs per node.

Yamazaki, Heinlein, Rajamanickam (2023)

Related works

- One-level Schwarz with local solves on GPUs: Luo, Yang, Zhao, Cai (2011)
- Solves of dense local Schur complement matrices in the balancing domain decomposition by constraints (BDDC) method on GPUs: Šístek & Oberhuber (2022)

Learning Extension Operators Using Graph Neural Networks

Why Learning Extension Operators

Most coarse spaces for Schwarz preconditioners are constructed based on a characteristic functions

$$\varphi_i(\omega_j)=\delta_{ij},$$

on specifically chosen sets of nodes $\{\omega_j\}_j$. The values in the remaining nodes are then obtained by extending the values into the adjacent subdomains. Examples:



Why Learning Extension Operators

Most coarse spaces for Schwarz preconditioners are constructed based on a characteristic functions

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Observation 1

Energy-minimizing extensions

are algebraic:

 $\mathbf{v}_{l} = -\mathbf{K}_{ll}^{-1}\mathbf{K}_{l\Gamma}\mathbf{v}_{\Gamma}$

(with Dirichlet b. c.)

 can be costly: solving a problem in the interior





Heterogeneous: $\alpha_{\text{light}} = 1$; $\alpha_{\text{dark}} = 10^8$

The performance may strongly depend on extension operator:

coarse space	its.	ĸ
—	163	$4.06 \cdot 10^7$
Q1	138	$1.07\cdot 10^{6}$
MsFEM	24	8.05

 \rightarrow Improving efficiency & robustness via machine learning.



- Lagrangian: geometric ext.
- MsFEM: geometric and energy-minimizing exts.
- RGDSW: algebraic and energy-minimizing exts.

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Related Works

This overview is not exhaustive:

Coarse spaces for domain decomposition methods

- Prediction of the geometric location of adaptive constraints (adaptive BDDC & FETI–DP as well as AGDSW): Heinlein, Klawonn, Lanser, Weber (2019, 2020, 2021, 2021, 2021, 2022)
- Prediction of the adaptive constraints: Klawonn, Lanser, Weber (preprint 2023, 2024)
- Prediction of spectral coarse spaces for BDDC for stochastic heterogeneities: Chung, Kim, Lam, Zhao (2021)
- Learning interface conditions and coarse interpolation operators: Taghibakhshi et al. (2022, 2023)

Algebraic multigrid (AMG)

- Prediction of coarse grid operators: Luz et al. (2020), Tomasi, Krause (2023)
- Coarsening: Taghibakhshi, MacLachlan, Olson, West (2021); Antonietti, Caldana, Dede (2023)

An overviews of the state-of-the-art on domain decomposition and machine learning in early 2021 and 2023:



A. Heinlein, A. Klawonn, M. Lanser, J. Weber

Combining machine learning and domain decomposition methods for the solution of partial differential equations — A review GAMM-Mitteilungen. 2021.



A. Klawonn, M. Lanser, J. Weber

Machine learning and domain decomposition methods – a survey arXiv:2312.14050...2023

Prediction via Graph Convolutional Networks

Graph neural networks (GNNs) introduced in Gori, Monfardini, and Scarselli (2005) are well-suited for learning on data based on simulation meshes:

- Generalization of classical convolutional neural networks (CNNs) LeCun (1998) to graph-based data sets.
- Aggregation and transmission of features of neighboring nodes in the graph via message passing layers.
- Invariance and equivariance with respect to position and permutation of the nodes of the graph.

Local approach

- Input: subdomain matrix K_i
- Output: basis functions {φ_j^{Ω_i}}
 on the same subdomain
- Training on subdomains with varying geometry
- Inference on unseen subdomains







Theory-Inspired Design of the GNN-Based Coarse Space

Null space property

Any extension-based coarse space built from a partition of unity on the domain decomposition interface satisfies the **null space property necessary for numerical scalability**:



Explicit partition of unity

To **explicitly enforce** that the basis functions $(\varphi_j)_i$ form a partition of unity

$$\varphi_j = \frac{\hat{\varphi}_j}{\sum_k \hat{\varphi}_k},$$

where the $\hat{\varphi}_k$ are the outputs of the GNN.

Initial and target

- Initial function: partition of unity that is constant in the interior
- Target function:
 - linear on the edges
 - energy-minimizing in the interior
- $\rightarrow \mbox{ Information transport via} \\ \mbox{ message passing }$





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Numerical Results for Homogeneous Laplacian – Weak Scaling Study

Model problem: 2D Laplacian model problem discretized using finite differences on a structured grid

$$-\Delta u = 1$$
 in Ω ,

$$u = 0$$
 on $\partial \Omega$,

decomposed using METIS:



 The GNN has been trained on 64 subdomains.



Yamazaki, Heinlein, Rajamanickam (in prep.)

Numerical Results for Heterogeneous Laplacian – Weak Scaling Study

Heterogeneous Laplacian with $\alpha_{max}/\alpha_{min} = 10^3$:



 $-\nabla \cdot (\alpha(x)\nabla u(x)) = f \text{ in } \Omega = [0,1]^2, \qquad u = 0 \text{ on } \partial\Omega.$

Yamazaki, Heinlein, Rajamanickam (in prep.)

FROSch

 FROSCH is based on the Schwarz framework and energy-minimizing coarse spaces, which provide numerical scalability using only algebraic information for a variety of applications

Subdomain solves on GPUs

- Subdomain solves make up a major part of the total solver time.
- Using the GPU triangular solve from KOKKOSKERNELS, we can speed up the solve phase of FROSCH. It can be further improved using ILU.

Learning extension operators

- Extensions are a major component in the construction of coarse spaces for domain decomposition methods.
- Using GNNs and known properties from the theory, we can learn extension operators that lead to a scalable coarse spaces.

Thank you for your attention!